

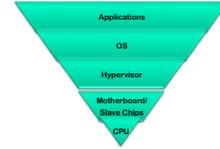
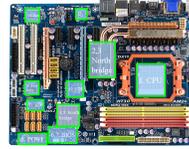
Tamper Evident Microprocessors

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Modern Hardware is Complex

- Modern systems built on layers of hardware



- Complexity increases risk of backdoors
 - More hands
 - Easier to hide
- A significant vulnerability
 - Hardware is the root of trust
 - All hardware and software controlled by microprocessors

Prior Work and Scope

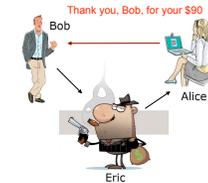
- Microprocessor design stages



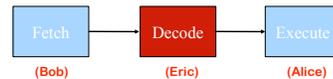
- Prior work focuses on back end
 - More immediate threat
 - Example: IC fingerprinting [Agrawal et al., 2007]
- Front end is the extreme root
 - Common assumption: golden model from front end
 - Focus of this work

Key Idea: Use Inherent Division of Work

- Bob
 - Nice Guy
 - Donates \$100
- Eric
 - Evil Accountant
 - Steals \$10
- Alice
 - Charity President
 - Receives \$90



Microprocessor Pipeline Stages Analogue

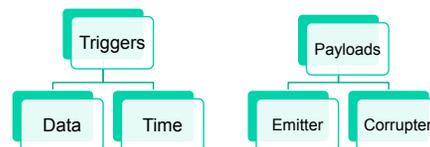


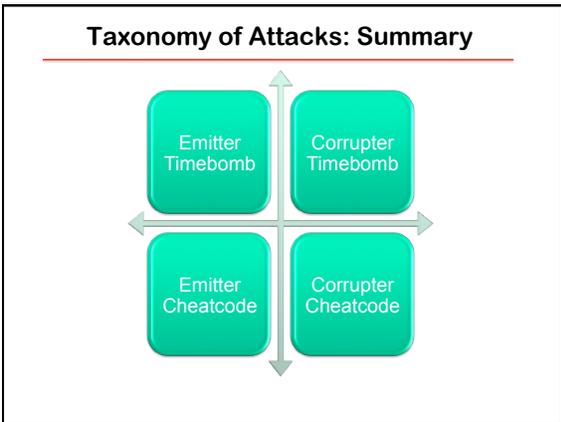
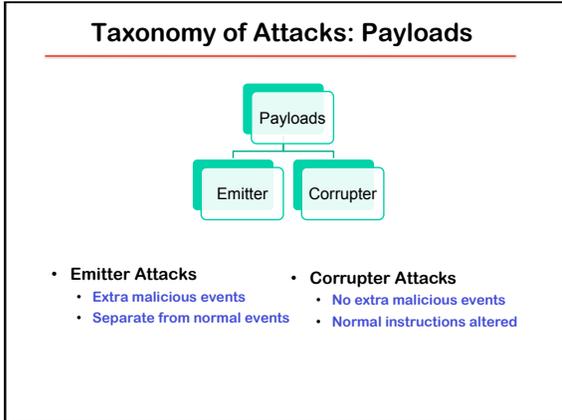
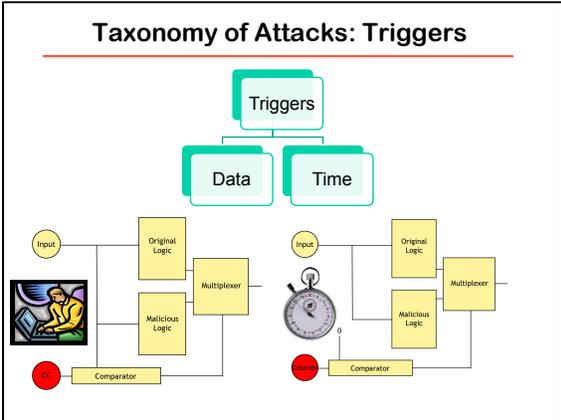
Outline

- Taxonomy
 - Ticking Timebombs, Cheat Codes, Emitters, Corrupters
- Solutions
 - TrustNet and DataWatch
- Results
 - Correctness, Coverage and Costs
- Future Work

Taxonomy of Attacks

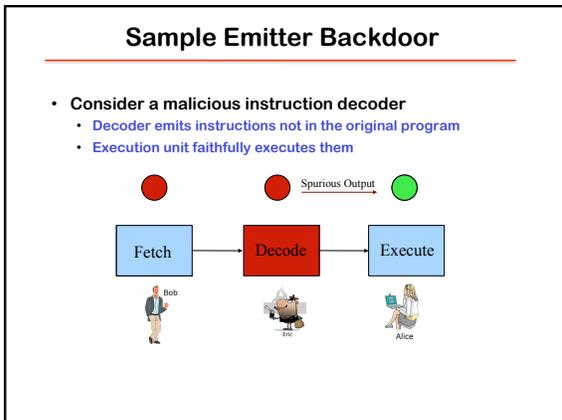
- Backdoor = Trigger + Payload
 - Trigger: Turns on an attack
 - Payload: Malicious, illegal action





- ### Assumptions
- **Large design team**
 - Each designer works on one unit or part of one
 - Security add-ons cannot be done by one member
 - **Full knowledge**
 - Attacker has complete access to all design specifications
 - Attacker also knows about additional security mechanism
 - **Equal distrust**
 - Any one designer/unit may be evil
 - Security add-ons may contain backdoors

- ### Outline
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TrustNet

- **Predictor and Reactor monitor the Target**
 - Division of work prevents *one* bad guy from breaking two units
 - Scaling to larger number increases design complexity

Corrupter Backdoors

- **Bob**
 - Still nice
 - Donates \$100
- **Eric**
 - Evil (and smarter)
 - Converts to Canadian \$
- **Alice**
 - Still president
 - Fooled by Eric's C\$100

DataWatch

- **Scaled up version of TrustNet**
 - Multiple bit messages
 - Confirms types of messages (instead of just yes/no)

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Experimental Context, Correctness, Costs

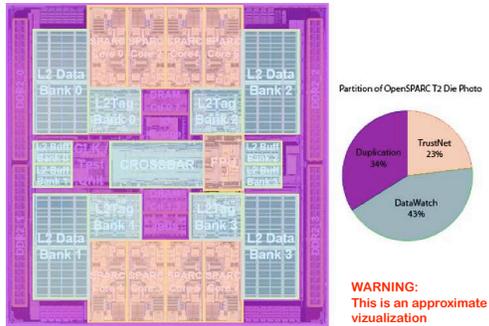
- **Context**
 - Simplified OpenSPARC T2
- **Correctness**
 - Designed attacks
 - No false positives or negatives
- **Costs**
 - Low area overhead (2 KB per core)
 - No performance impact
- **How to measure coverage?**

Coverage: Vulnerability Space

Units with a core Units with a core

Paper has plots for other units at a chip level

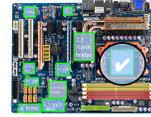
Coverage Visualization



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Summary and Future Work

- **Strengthen root of trust: microprocessors**
 - Hardware-only solution. No perf impact, low area overhead
 - Security add-on highly resilient to corruption
 - Provided attack taxonomy, method to characterize attack space
- **Applicability of TrustNet & DataWatch**
 - Covered: pipelines, caches and content associative memory
 - Not covered: ALU, microcode, power mgmt., side-channels
- **Moving Forward**
 - Expand coverage
 - Out-of-order processors
 - Motherboard components
 - Design automation tools
 - Reaction to errors
 - Applying techniques for reliable execution
- First steps toward a secure trusted hardware w/ untrusted units



Thank You! and Questions?